

# Minimal interpolants in $C^{1,1}(\mathbb{R}^n)$

Matthew J. Hirn

Department of Mathematics  
Yale University

École normale supérieur de Cachan  
Groupe de travail “application des mathématiques”

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# Overview

- 1 Introduction
- 2 A minimal interpolant
  - An interpolant (Wells)
  - The minimal norm (Le Gruyer)
  - Le Gruyer + Wells
- 3 Wells' construction
  - Dual cell complexes
  - Defining the interpolant
- 4 Implementation on a computer
  - The model of computation
  - Efficient computation for  $C^{1,1}(\mathbb{R}^2)$

## Our setting

- Define the Lipschitz semi-norm of the gradient of a function  $F : \mathbb{R}^n \rightarrow \mathbb{R}$  as:

$$\text{Lip}(\nabla F) \triangleq \sup_{\substack{x, y \in \mathbb{R}^n \\ x \neq y}} \frac{\|\nabla F(x) - \nabla F(y)\|}{\|x - y\|}$$

- We will work in  $C^{1,1}(\mathbb{R}^n)$ , which is defined as:

$$C^{1,1}(\mathbb{R}^n) \triangleq \{F \in C^1(\mathbb{R}^n) \mid \text{Lip}(\nabla F) < \infty\}$$

# Input data

We are given the following data:

- The **set of points** that we wish to interpolate through. Denote this set as

$$E = \{p_1, \dots, p_N\} \subset \mathbb{R}^n$$

- The **function values** at each point of  $E$ :

$$\begin{aligned} f : E &\rightarrow \mathbb{R} \\ p &\mapsto f(p) \end{aligned}$$

- The **gradients** at each point of  $E$ :

$$\begin{aligned} f_{\nabla} : E &\rightarrow \mathbb{R}^n \\ p &\mapsto f_{\nabla}(p) \end{aligned}$$

# Minimal interpolant

Given  $(E, f, f_{\nabla})$ , we want to compute a minimal interpolant  $F \in C^{1,1}(\mathbb{R}^n)$  such that:

- **Interpolate:**

$$\begin{aligned} F(p) &= f(p), & \text{for all } p \in E \\ \nabla F(p) &= f_{\nabla}(p), & \text{for all } p \in E \end{aligned}$$

- **Minimum norm:** If  $\tilde{F} \in C^{1,1}(\mathbb{R}^n)$  is another interpolant for the data  $(E, f, f_{\nabla})$ , then

$$\text{Lip}(\nabla F) \leq \text{Lip}(\nabla \tilde{F})$$

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# Wells' theorem

- Let  $M > 0$ , and suppose for now that we just want an interpolant  $F$  for the data  $(E, f, f_{\nabla})$  such that  $\text{Lip}(\nabla F) = M$ . When is this possible?

## Theorem (Wells, 1973)

Let  $M > 0$  and suppose you are given data  $(E, f, f_{\nabla})$ . If for each pair  $(p, p') \in E \times E$ ,

$$f(p') \leq f(p) + \frac{1}{2} \langle f_{\nabla}(p) + f_{\nabla}(p'), p' - p \rangle + \frac{M}{4} \|p' - p\|^2 - \frac{1}{4M} \|f_{\nabla}(p') - f_{\nabla}(p)\|^2,$$

then there exists an interpolant  $F \in C^{1,1}(\mathbb{R}^n)$  for the data  $(E, f, f_{\nabla})$  such that  $\text{Lip}(\nabla F) = M$ . Furthermore you can explicitly construct such an  $F$ .

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# Le Gruyer's $\Gamma^1$ constant

- Define  $\text{Lip}(E, f, f_{\nabla})$  as:

$$\text{Lip}(E, f, f_{\nabla}) \triangleq \inf\{\text{Lip}(\nabla F) \mid F \text{ interpolates the data } (E, f, f_{\nabla})\}$$

## Definition (Le Gruyer, 2009)

Also define  $\Gamma^1(E, f, f_{\nabla})$  as

$$\Gamma^1(E, f, f_{\nabla}) \triangleq 2 \max_{\substack{p, p' \in E \\ p \neq p'}} \sup_{x \in \mathbb{R}^n} \frac{T_p(x) - T_{p'}(x)}{\|p - x\|^2 + \|p' - x\|^2},$$

where  $T_p(x) \triangleq f(p) + \langle f_{\nabla}(p), x - p \rangle$ .

## Theorem (Le Gruyer, 2009)

For any data  $(E, f, f_{\nabla})$ ,

$$\Gamma^1(E, f, f_{\nabla}) = \text{Lip}(E, f, f_{\nabla})$$

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# Combining the results of Le Gruyer and Wells

## Mini-lemma

Given data  $(E, f, f_{\nabla})$ , Le Gruyer's constant  $\Gamma^1 = \Gamma^1(E, f, f_{\nabla})$  satisfies the condition of Wells, i.e., for each pair  $(p, p') \in E \times E$ ,

$$f(p') \leq f(p) + \frac{1}{2} \langle f_{\nabla}(p) + f_{\nabla}(p'), p' - p \rangle + \frac{\Gamma^1}{4} \|p' - p\|^2 - \frac{1}{4\Gamma^1} \|f_{\nabla}(p') - f_{\nabla}(p)\|^2$$

## Theorem (Le Gruyer & Wells)

Given data  $(E, f, f_{\nabla})$ , one can construct a minimal interpolant  $F \in C^{1,1}(\mathbb{R}^n)$  for  $(E, f, f_{\nabla})$ .

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# Cell complex

## Definition

A **cell complex** in  $\mathbb{R}^n$  is a set  $\mathcal{K}$  of convex polyhedra (called cells) in  $\mathbb{R}^n$  satisfying two conditions:

- 1 Every face of a cell is also a cell in  $\mathcal{K}$
- 2 If  $\sigma_1, \sigma_2 \in \mathcal{K}$  and  $\sigma_1 \cap \sigma_2 \neq \emptyset$ , then  $\sigma_1 \cap \sigma_2$  is a face of both  $\sigma_1$  and  $\sigma_2$

# Shift the points in $E$

- For each  $p \in E$ , define  $\tilde{p} \in \mathbb{R}^n$  as

$$\tilde{p} \triangleq p - \frac{f_{\nabla}(p)}{\Gamma^1}$$

- Furthermore, for each  $S \subseteq E$ , set

$$\tilde{S} \triangleq \{\tilde{p} \mid p \in S\}$$

# Some useful functions

- For each  $p \in E$ , define  $d_p : \mathbb{R}^n \rightarrow \mathbb{R}$  as

$$d_p(x) \triangleq f(p) - \frac{1}{2\Gamma^1} \|f_{\nabla}(p)\|^2 + \frac{\Gamma^1}{4} \|x - \tilde{p}\|^2$$

- Furthermore, for each  $S \subseteq E$ , define  $d_S : \mathbb{R}^n \rightarrow \mathbb{R}$  as

$$d_S(x) \triangleq \min_{p \in S} d_p(x)$$

# Special subsets of $E$

- For each  $S \subseteq E$ , define

$$S_e \triangleq \{x \in \mathbb{R}^n \mid d_p(x) = d_{p'}(x), \forall p, p' \in S\}$$

- Also for each  $S \subseteq E$ , define

$$S_* \triangleq \{x \in \mathbb{R}^n \mid d_p(x) = d_{p'}(x) \leq d_{p''}(x), \forall p, p' \in S, \forall p'' \in E\}$$

## Definition

$$K \triangleq \{S \subseteq E \mid \text{for some } x \in S_*, d_S(x) < d_{E \setminus S}(x)\}$$



# The cell complex $\mathcal{K}_*$

## Definition

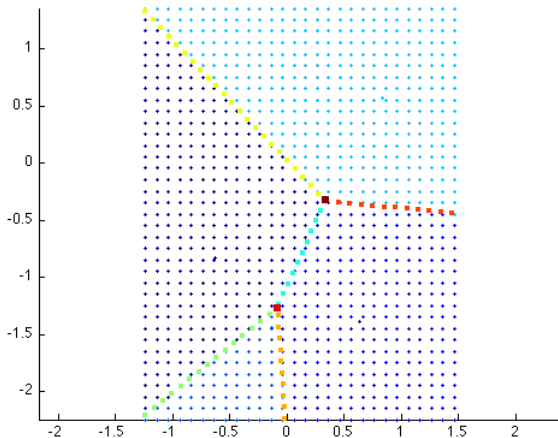
$$\mathcal{K}_* \triangleq \{S_* \mid S \in K\}$$

## Lemma (Wells, 1973)

$\mathcal{K}_*$  is a cell complex in  $\mathbb{R}^n$ , and furthermore

$$\bigcup_{S \in K} S_* = \mathbb{R}^n$$

# The cell complex $\mathcal{K}_*$



# The cell complex $\widehat{\mathcal{K}}$

- For each  $S \subseteq E$ , define

$S_a \triangleq$  the smallest affine subspace of  $\mathbb{R}^n$  containing  $\widetilde{S}$

- Also for each  $S \subseteq E$ , define

$\widehat{S} \triangleq$  convex hull of  $\widetilde{S}$

## Definition

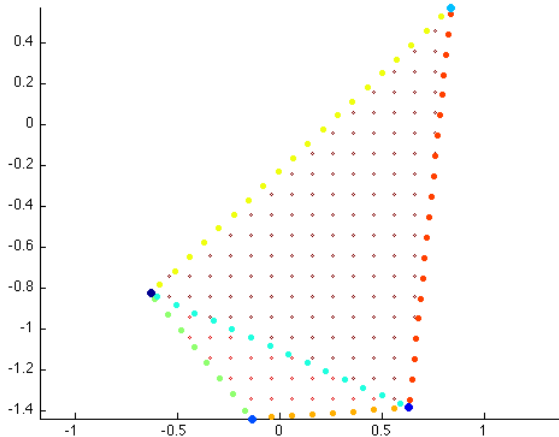
$$\widehat{\mathcal{K}} \triangleq \{\widehat{S} \mid S \in K\}$$

## Lemma (Wells, 1973)

$\widehat{\mathcal{K}}$  is a cell complex in  $\mathbb{R}^n$ , and furthermore

$$\bigcup_{S \in K} \widehat{S} = \text{convex hull of } \widetilde{E}$$

# The cell complex $\widehat{\mathcal{K}}$



# Duality between $\widehat{\mathcal{K}}$ and $\mathcal{K}_*$

The two cell complexes  $\widehat{\mathcal{K}}$  and  $\mathcal{K}_*$  have the following **duality** between them:

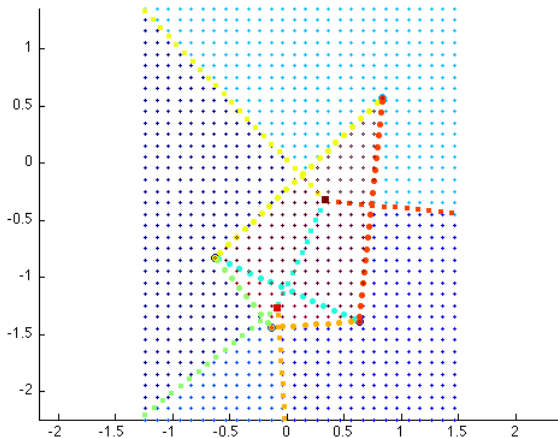
- For each  $S \in \mathcal{K}$ ,

$$\widehat{S} \perp S_* \quad \text{and} \quad \dim(\widehat{S}) + \dim(S_*) = n$$

- If  $S, S' \in \mathcal{K}$  and  $S \subseteq S'$ , then

$$\widehat{S} \subseteq \widehat{S'} \quad \text{and} \quad S'_* \subseteq S_*$$

# The cell complexes $\widehat{\mathcal{K}}$ and $\mathcal{K}_*$



# A third cell complex

## Definition

For each  $S \in K$ , let

$$T_S \triangleq \{x \in \mathbb{R}^n \mid x = \frac{1}{2}(y + z) \text{ for some } y \in \widehat{S} \text{ and } z \in S_*\},$$

and set

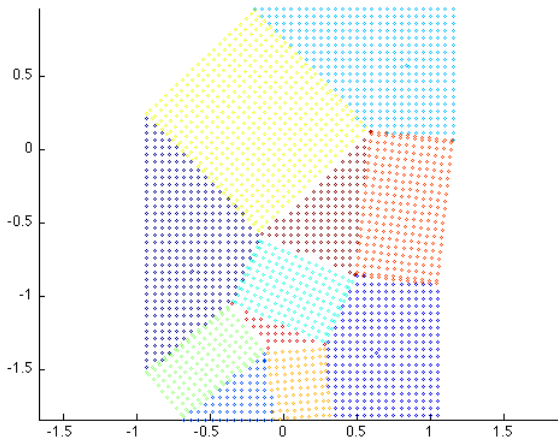
$$\mathcal{K}_T \triangleq \{T_S \mid S \in K\}$$

## Lemma (Wells, 1973)

$\mathcal{K}_T$  is a cell complex in  $\mathbb{R}^n$ , and furthermore

$$\bigcup_{S \in K} T_S = \mathbb{R}^n$$

# The cell complex $\mathcal{K}_T$





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# The minimal interpolant

- For each  $S \in K$ , set  $S_c \triangleq S_a \cap S_e \in \mathbb{R}^n$  (note:  $S_a \perp S_e$ ).

## Definition

For each  $S \in K$ , define  $F_S : T_S \rightarrow \mathbb{R}$  as

$$F_S(x) \triangleq d_S(S_c) + \frac{\Gamma^1}{2} (\text{dist}(x, S_a)^2 - \text{dist}(x, S_e)^2), \quad \forall x \in T_S$$

## Theorem (Wells and Le Gruyer)

A minimal interpolant  $F : \mathbb{R}^n \rightarrow \mathbb{R}$  can then be defined as:

$$F(x) \triangleq F_S(x), \quad \text{if } x \in T_S$$

# Proof of theorem

## Lemma (Wells, 1973)

For each  $S \in K$ ,  $F_S \in C^\infty(T_S)$  and furthermore,

$$\nabla F_S(x) = \frac{\Gamma^1}{2}(z - y),$$

where  $x = \frac{1}{2}(y + z)$  and  $y \in \widehat{S}$ ,  $z \in S_*$ .

# Proof of theorem

Now one needs to check several things:

- Is  $F$  well defined?

$$F_S(x) = F_{S'}(x), \text{ for all } x \in T_S \cap T_{S'}$$

Proof: use  $S_a \perp S_e, \widehat{S} \perp S_*$ , and  $S_c = S_a \cap S_e$

- Is  $F \in C^{1,1}(\mathbb{R}^n)$ ?

$$\nabla F_S(x) = \nabla F_{S'}(x), \text{ for all } x \in T_S \cap T_{S'}$$

Proof: use the previous lemma

- Does  $\text{Lip}(\nabla F) = \Gamma^1$ ?

$$\text{Lip}(\nabla F_S) = \Gamma^1$$

Proof: use previous lemma and  $\widehat{S} \perp S_*$

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# Proof of theorem

Finally, we need to check that  $F$  interpolates the data  $(E, f, f_{\nabla})$ :

$$F(p) = f(p) \quad \text{and} \quad \nabla F(p) = f_{\nabla}(p), \quad \forall p \in E$$

Proof (fix  $p \in E$ ):

- Clearly,  $\{p\}_a = \widehat{\{p\}} = \{\bar{p}\} = \{p - \frac{f_{\nabla}(p)}{\Gamma^1}\} \Rightarrow \{p\}_c = \bar{p}$
- Wells' condition  $\Rightarrow d_p(p + \frac{f_{\nabla}(p)}{\Gamma^1}) \leq d_{p'}(p + \frac{f_{\nabla}(p)}{\Gamma^1}), \quad \forall p' \in E$
- $\Rightarrow p + \frac{f_{\nabla}(p)}{\Gamma^1} \in \{p\}_*$
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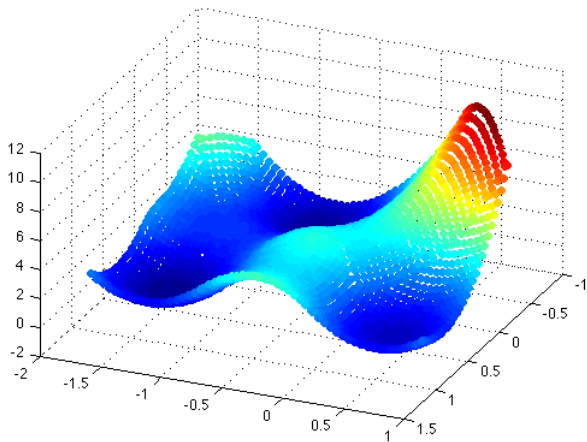
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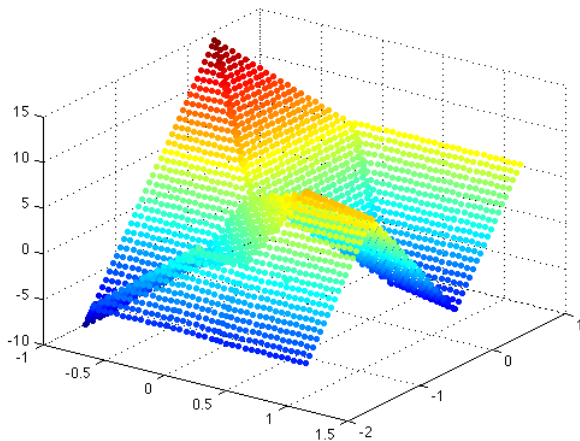
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# The minimal interpolant $F$





# The partial derivative $\frac{\partial F}{\partial y}$



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  - The minimal norm (Le Gruyer)
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# Computer model

We want to **efficiently compute** minimal interpolants  $F \in C^{1,1}(\mathbb{R}^n)$  for the data  $(E, f, f_{\nabla})$ .

Our computer model is the following:

- It can work with real numbers and an exact real number can be stored at each memory address
- It takes one machine operation to add, subtract, multiply, or divide two real numbers  $x$  and  $y$ , or to compare them (i.e., decide whether  $x < y$ ,  $x > y$ , or  $x = y$ )

# Compute a function $F$

To “compute a function  $F$ ” means the following:

- First, we enter the data  $(E, f, f_{\nabla})$  into the computer.
- The computer then works for a while, performing  $L_0$  machine operations. This is the **one-time work**.
- The computer then signals that it is ready to accept further input.
- Whenever we enter a point  $x \in \mathbb{R}^n$ , the computer responds by producing  $(F(x), \nabla F(x))$ , using  $L_1$  machine operations to make the computation. This is the **query work**.
- The **storage** of our algorithm is the number of memory addresses required to carry out the above work.

# Efficiency goals

Recall,  $|E| = N$ .

Ideally, our algorithm will only require the following:

- One-time work:  $C(n)N \log N$
- Query work:  $C(n) \log N$
- Storage:  $C(n)N$

Here,  $C(n)$  is a constant depending only on  $n$ .

# Overview

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# Computation of $\Gamma^1$

- To compute a minimal interpolant  $F$  *exactly*, using only  $C(n)N \log N$  one-time work, is most likely impossible.
- Indeed, recall the constant  $\Gamma^1$ , defined as:

$$\Gamma^1 \triangleq 2 \max_{\substack{p, p' \in E \\ p \neq p'}} \sup_{x \in \mathbb{R}^n} \frac{T_p(x) - T_{p'}(x)}{\|p - x\|^2 + \|p' - x\|^2}$$

- In fact, by a result of Le Gruyer (2009), one can write  $\Gamma^1$  as:

$$\Gamma^1 = \max_{\substack{p, p' \in E \\ p \neq p'}} \sqrt{A_{p,p'}^2 + B_{p,p'}^2} + |A_{p,p'}|$$

- Still though, to compute  $\Gamma^1$ , we must use  $C(n)N^2$  operations.

# Approximation of $\Gamma^1$

- We propose computing an approximation  $\tilde{\Gamma}^1$  that is within  $\varepsilon > 0$  of  $\Gamma^1$ , i.e.,

$$\tilde{\Gamma}^1 \leq (1 + \varepsilon)\Gamma^1$$

## Conjecture/Lemma (H. and Narayanan, 2011)

Let  $\varepsilon > 0$  and suppose you are given data  $(E, f, f_\nabla)$ , where  $|E| = N$ . Then there exists an algorithm, using no more than  $C(n, \varepsilon)N$  storage and  $C(n, \varepsilon)N \log N$  work, that computes a constant  $\tilde{\Gamma}^1(E, f, f_\nabla)$  such that

$$\tilde{\Gamma}^1 \leq (1 + \varepsilon)\Gamma^1$$

- Main idea: use a well separated pairs decomposition (WSPD) for  $E$  in the same spirit as how one approximates  $\text{Lip}(f)$ . By Callahan and Kosaraju (1995), the WSPD can be computed using only  $C(n, \varepsilon)N$  storage and  $C(n, \varepsilon)N \log N$  work.



# Computation of $F$

## Conjecture/Theorem (H. and Narayanan, 2011)

Let  $\varepsilon > 0$  and suppose you are given data  $(E, f, f_{\nabla})$ , where  $E \subset \mathbb{R}^2$  and  $|E| = N$ . Then there exists an algorithm, using no more than  $C(\varepsilon)N$  storage,  $C(\varepsilon)N \log N$  one-time work, and  $C(\varepsilon) \log N$  query work, that computes an interpolant  $F$  for the data  $(E, f, f_{\nabla})$  such that  $\text{Lip}(\nabla F) \leq (1 + \varepsilon)\Gamma^1$ .

The main ideas are the following:

- Use the previous lemma to compute  $\tilde{\Gamma}^1$  (one-time work).
- The cell complex  $\mathcal{K}_*$  is a generalization of a Voronoi diagram, which in  $\mathbb{R}^2$  can be computed using  $CN$  storage and  $CN \log N$  operations (one-time work).
- Given  $x \in \mathbb{R}^2$ , determining which  $T_S$  cell contains  $x$  is a point location problem. In  $\mathbb{R}^2$ , this can be done using  $CN$  storage and  $\log N$  operations (query work).

# Future directions

Some possible future directions are:

- Computation of interpolants in higher dimensions.
- Consider the same space  $C^{1,1}(\mathbb{R}^n)$ , but suppose you are only given data  $(E, f)$  (i.e. no information on the gradients).
- Consider higher order derivatives, i.e.,  $C^{m,1}(\mathbb{R}^n)$ .

Merci!

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